

LAST TIME: Given  $\mathcal{L}^\otimes$  symm. mon.  $\omega$ -cat. a commutative algebra in  $\mathcal{L}^\otimes$  is

$$A: *^{\text{Fin}} \rightarrow \mathcal{L}^{\otimes, \text{Fin}} \quad \text{which sends}$$

inert  $\omega$ -Cartesian morphisms to inert  $\omega$ -Cartesian morphisms.

Explicitly, let  $e_i: \langle 2 \rangle \rightarrow \langle 1 \rangle$  be  $e_i(j) = \delta_{ij}$  two inert morphisms. Then  $A(e_i): A\langle 2 \rangle \rightarrow A\langle 1 \rangle$  is  $\omega$ -Cartesian if  $\forall z \in \mathcal{L}^{\otimes, \text{Fin}}$  one has:

$$(\Delta) \quad \text{Hom}_{\mathcal{L}^{\otimes, \text{Fin}}} (A\langle 1 \rangle, z) \xrightarrow{\cong} \text{Hom}_{\mathcal{L}^{\otimes, \text{Fin}}} (A\langle 2 \rangle, z) \times_{\text{Hom}_{\text{Fin}} (\langle 2 \rangle, p(z))} \text{Hom}_{\text{Fin}} (\langle 1 \rangle, z).$$

Recall that a morphism  $A\langle 1 \rangle \rightarrow z$  in  $\mathcal{L}^{\otimes, \text{Fin}}$  is the data of:

$$f: \langle 1 \rangle \rightarrow p(z) \quad + \quad \bar{f}: \mathcal{L}^\otimes(f)(A\langle 1 \rangle) \rightarrow \bar{z},$$

where  $\bar{z} \in \mathcal{L}^\otimes(p(z))$ .

Similarly, let  $g: \langle 2 \rangle \rightarrow p(z) \quad + \quad \bar{g}: \mathcal{L}^\otimes(g)(A\langle 2 \rangle) \rightarrow \bar{z}$  be an elab on the right-hand side.

The condition that  $(\Delta)$  is an isom. gives:

$$g = f \circ e_i \quad \text{and} \quad \bar{g} = \bar{f} \circ \mathcal{L}^\otimes(e_i).$$

Since we require  $(\Delta)$  holds for all  $z \in \mathcal{L}^{\otimes, \text{Fin}}$ , by considering  $\bar{z} \in \mathcal{L}^\otimes(\langle 1 \rangle)$  and  $f = \text{id}_{\langle 1 \rangle}$  one gets.

$$\text{Hom}_{\mathcal{L}^\otimes} (A\langle 1 \rangle, \bar{z}) = \text{Hom}_{\mathcal{L}^\otimes} (\mathcal{L}^\otimes(e_i)(A\langle 2 \rangle), \bar{z})$$

$\forall \bar{z} \in \mathcal{L}^\otimes$ .

$$\Rightarrow \mathcal{L}^\otimes(e_i)(A\langle 2 \rangle) = A\langle 1 \rangle.$$

Thus, by applying the description of  $\otimes$  to the image of  $A(\alpha)$  we get:

$$A(\alpha) =: A\langle 1 \rangle \otimes A\langle 1 \rangle \rightarrow A\langle 1 \rangle.$$

However, notice that ~~we~~ requiring that  $A(\alpha)$  is  $\otimes$  cat. would imply  $\mathcal{L}^{\otimes(\alpha)}(A\langle 2 \rangle) \stackrel{!}{=} A\langle 1 \rangle$  in  $\mathcal{L}$ , which is  $A\langle 1 \rangle \otimes A\langle 1 \rangle$ , which

we don't want to in general.

(Discussion of monoidal categories see last time. Maybe postponed.)

Here is another important example of a sym. mon.  $\infty$ -category, which is a sanity check for our definition.

Example: (ord. ~~mon~~ <sup>tensor</sup>  $\otimes$ -~~mon~~ cat.) let  $C$  be a sym. mon. 1-cat. Consider the coCartesian fibration  $C^{\otimes, \text{Fin}} \rightarrow \text{Fin}$  defined as follows:

-  $C^{\otimes, \text{Fin}}_{\langle n \rangle} = \{ \{ X_i \}_{i \in \langle n \rangle} \mid X_i \in C \}$   
a morphism

-  $(\{ X_i \}_{i \in \langle n \rangle}) \rightarrow (\{ Y_j \}_{j \in \langle m \rangle})$  is:

$$\alpha: \langle n \rangle \rightarrow \langle m \rangle + \tilde{\alpha} \in \prod_{j \in \langle m \rangle} \text{Hom}_C \left( \bigotimes_{i \in \alpha^{-1}(j)} X_i, Y_j \right)$$

where  $X_0 := 1_C$ .

Exercise: Check the above gives a sym. mon.  $\infty$ -category.

Sanity check 2: For  $C^{\otimes, \text{Fin}} \rightarrow \text{Fin}_*$  a sym. mon.  $\infty$ -cat. obtained as in the previous example.

$A \in \text{CAlg}(C^\bullet)$  a comm. alg. object in the sense of last time.

is equivalent to  $\bar{A} \in C$ ,  $+ \mu: \bar{A} \otimes \bar{A} \rightarrow \bar{A}$  a commutative and associative multiplication.

Indeed, the remark from earlier today gives  $A\langle 2 \rangle = A\langle 1 \rangle \otimes A\langle 1 \rangle$ .  
So  $\bar{A} = A\langle 1 \rangle$ .

Associativity is the commutativity of the ~~Rees~~ image of

$$\begin{array}{ccc} \begin{pmatrix} 1 & 2, 3 \\ \downarrow & \downarrow \\ 1 & 2 \end{pmatrix} & \langle 3 \rangle & \xrightarrow{(1, 2 \rightarrow 1, 3 \rightarrow 2)} \langle 2 \rangle \\ \downarrow & \downarrow & \downarrow \alpha \\ 1 & 2 & \langle 2 \rangle \xrightarrow{\alpha} \langle 1 \rangle \end{array}$$

Similarly for commutativity.

That one can write  $A: \text{Fin}_* \rightarrow \text{Cat} \rightarrow \text{Set}$  for  $A$  is an easy check from the description of  $C^{\otimes, \text{Fin}}$ .

To compare the notions of derived rings to usual rings we need to quickly discuss one <sup>more</sup> <sup>about</sup>  $\infty$ -categories.

Def'n: let  $X \in \mathcal{C}$  be an object in an  $\infty$ -cat.  $X$  is said to be  $n$ -truncated ( $n \geq 0$ ) if  $\forall A \in \mathcal{C}$  we have:  
 $\text{Hom}_{\mathcal{C}}(X, A) \in \text{Spc}^{\leq n} = \{ \mathbb{Q} \otimes S \in \text{Spc} \mid \hat{n}(S) = 0 \}$   
 $\forall i > n$

Let  $\mathcal{Z}_{\leq n}^{\otimes} \mathcal{C}$  denote the <sup>sub</sup>category of  $n$ -truncated objects.  
In particular, objects of  $\mathcal{Z}_{\leq 0}^{\otimes} \mathcal{C}$  are called discrete objects.

Here are a couple of useful properties:

- Lemma:
- (i)  $\forall n \geq 0, \sum_{\leq n} \mathcal{L} \hookrightarrow \mathcal{L}$  is stable under limits.
  - (ii)  ~~$\mathcal{L}$~~   $\sum_{\leq n} \mathcal{L} \rightarrow \mathcal{L}$  has a left adjoint. (that exist in  $\mathcal{L}$ )
  - (iii) any functor that preserves finite limits will preserve  $n$ -truncated objects.

- Example:
- (i)  $\sum_{\leq 0} \text{Spc} = \text{sets}$ .
  - (ii)  $\sum_{\leq 1} \text{Spc} = \text{Grpd}$ .

To give more examples, we recall the following notion:

Def'n: (Chronologically graded t-structure) Let  $D$  be a triangulated category. a t-structure on  $D$  is the data of two subcategories.  $(D^{\geq 0}, D^{\leq 0})$  s.t.:

(i)  $D^{\leq -1} := [1](D^{\leq 0}) \subseteq D^{\leq 0}$  and  $D^{\geq 1} := [-1](D^{\geq 0}) \subseteq D^{\geq 0}$ ;

(ii)  $\forall X \in D^{\leq 0}$  and  $Y \in D^{\geq 1}$   
 $\text{Hom}_D(X, Y) = 0$ .

(iii)  $\forall X \in D, \exists X_{\text{con}} \in X^{\leq 0}$ , and  $X_{\text{cocon}} \in X^{\geq 1}$   
 and a distinguished triangle:

$$X_{\text{con}} \rightarrow X \rightarrow X_{\text{cocon}}$$

For  $\mathcal{L}$  a stable  $\omega$ -cat. a t-str. is the data of  $(\mathcal{L}^{\leq 0}, \mathcal{L}^{\geq 0})$  subcategories s.t.  $h\mathcal{L} = (h\mathcal{L}^{\leq 0}, h\mathcal{L}^{\geq 0})$  is a t-structure.

We ~~list~~ give a small and incomplete list of properties:

- Prop: (i)  $\mathcal{L}^{\heartsuit} := \mathcal{L}^{\geq 0} \cap \mathcal{L}^{\leq 0}$  is a 1-category and it is abelian;  
 (ii)  $\mathcal{L}^{\leq n} \hookrightarrow \mathcal{L}$  has a right adjoint  $\tau^{\leq n}$ ,  $\mathcal{L}^{\leq n}$  is stable under colimits;  
 (iii)  $\mathcal{L}^{\geq n} \hookrightarrow \mathcal{L}$  has a left adjoint  $\tau^{\geq n}$ ,  $\mathcal{L}^{\geq n}$  is stable under limits.

The notations  $\tau^{\leq n}$  and  $\tau^{\geq -n}$  agree only on the subcategory  $\mathcal{L}^{\leq 0}$ .

Claim:  $X \in \mathcal{L}^{\leq 0}$  is  $n$ -truncated.  $\iff \text{Hom}_{\mathcal{L}}(Y, X) \in \text{Spec}^{\leq k} \forall Y \in \mathcal{L}^{\leq 0}$ .  
 $\iff X \in \mathcal{L}^{\geq -k}$ .

Notice  $\text{h}_i(\text{Hom}_{\mathcal{L}}(Y, X)) = \text{h}_0(\text{Hom}_{\mathcal{L}}(Y, X[i])) = \text{Hom}_{\mathcal{L}}(Y[-i], X)$ .

$\text{Ext}_{\mathcal{L}}^{-i}(Y, X) = 0 \quad \forall Y \in \mathcal{L}^{\leq 0} \iff \text{Ext}^i(Y, X) = 0$

$\text{Hom}_{\mathcal{L}}(Y, X[-i]) = 0 \quad \forall Y \in \mathcal{L}^{\leq 0}$

$X[-i] \in \mathcal{L}^{\geq 1} \implies X \in \mathcal{L}^{\geq 1-i}$ , taking  $i = k+1 \implies X \in \mathcal{L}^{\geq -k}$ .

Thus, Example: (iii) For  $\mathcal{L}$  stable w/ a t-structure:  
 $\tau^{\leq n}(\mathcal{L}^{\leq 0}) \simeq \mathcal{L}^{\geq -n, \leq 0} := \mathcal{L}^{\geq -n} \cap \mathcal{L}^{\leq 0}$ .

Property of a t-str. on a triangulated category.  
 $X \in \mathcal{L}^{\geq 1}$   
 $\iff \text{Hom}_{\mathcal{L}}(Y, X) = 0 \forall Y \in \mathcal{L}^{\leq 0}$ .

As much as I am trying to avoid these notes have some circularity.  
 For ease of exposition now we will assume a result, which we might

"justify" later. ~~There~~

FACT: The  $\infty$ -category  $\text{Vect}_k := \mathcal{D}(k)$  has a sym. mon. structure (k will be a field for us, but this is true for any comm. ring.).

It is characterized by: (i)  $\otimes$  commutes w/ colimits in each variable. (ii) when restricted to  $\text{Vect}_k^{\heartsuit} =$  ordinary category of k-vector spaces: it agrees w/ the usual  $\otimes$ -product.

( $\text{Vect}_k$  has a t-structure:  $\text{Vect}_k^{\leq 0} := \{ V \in \text{Vect}_k \mid H^i(V) = 0 \forall i \geq 1 \}$   
 $\text{Vect}_k^{\geq 0} := \{ \quad \quad \quad \mid H^i(V) = 0 \forall i \leq -1 \}$ )

Notice that by Ex: (i:ii) above one has:

$$\mathcal{Z}_{\leq 0}(\text{Vect}_k^{\leq 0}) = \text{Vect}_k^{\geq 0, \leq 0} = \text{Vect}^{\heartsuit} \text{ which is the ordinary cat. of } k\text{-vector spaces.}$$

Moreover, b/c  $\text{Vect}^{\leq 0} \hookrightarrow \text{Vect}$  preserves colimits, the  $\otimes$ -str. of  $\text{Vect}$  restricts to a  $\otimes$ -str. on  $\text{Vect}^{\leq 0}$ .

Finally, we can define the first candidate for derived  $\mathbb{P}$ -rings. (or k-algebras).

Def'n: A connective differential graded k-algebra (cdga) is an object of:

$$\text{CAlg}_k := \text{CAlg}(\text{Vect}_k^{\leq 0}), \text{ or } \text{CAlg} := \text{Fun}^{\text{right-ex}}(\mathcal{L}^{\text{Fine}}, \text{Vect}_k^{\leq 0, \otimes, \text{Fine}}).$$

(we never formally wrote down what the category  $\text{CAlg}(\mathbb{Z})$  for some sym. mon.  $\infty$ -cat.  $\mathcal{L}$  is. the right-hand side above hints at what formal symbols can do this for us.)

Rk: Here is some justification for this definition.

First, we claim that  $\tau^{z-n} : \text{Vect}^{\leq 0} \rightarrow \text{Vect}^{z-n, \leq 0}$  is a sym. monoidal functor.

Heuristically, notice that for any  $V \xrightarrow{f} W$  in  $\text{Vect}^{\leq 0}$ , s.t.

$\tau^{z-n}(f)$  is an isomorphism one has:

$$\tau^{z-n}(V \otimes U) \cong \tau^{z-n}(W \otimes U) \text{ for}$$

any  $U \in \text{Vect}^{\leq 0}$ .

By applying this twice to  $V \rightarrow \tau^{z-n}(V)$  and  $U \rightarrow \tau^{z-n}(U)$  one gets:

$$\tau^{z-n}(V \otimes U) \cong \tau^{z-n}(V) \otimes \tau^{z-n}(U).$$

Lemma:  $\text{Vect}^{z-n, \leq 0} \hookrightarrow \text{Vect}^{\leq 0}$  is a right-lax functor.

This follows from the following general fact justifying the names. right/left lax functors.

Thus, given a comm. alg. object in  $\text{Vect}^{z-n, \leq 0}$ , i.e.  $\mathbb{Q}$  a classical  $\mathbb{Q}$  commutative  $k$ -algebra one obtains an object in  $\text{CAlg}_k$  via:

$$z : \text{CAlg}(\text{Vect}^{z-n, \leq 0}) \rightarrow \text{CAlg}(\text{Vect}^{\leq 0}) = \text{CAlg}_k$$

Moreover, we claim:

Prop:  $z$  is fully faithful  $\mathbb{Q}$  with essential image the discrete objects of  $\text{CAlg}(\text{Vect}^{\leq 0}) = \text{CAlg}_k$ .

Idea: We haven't introduced these objects but given a free alg. functor:

$\text{Sym} : \text{Vect}^{\leq 0} \rightarrow \text{CAlg}_k$ , this follows fr. the adjunction  $\text{Hom}_{\text{CAlg}_k}(\text{Sym}(V), A) = \text{Hom}_{\text{Vect}^{\leq 0}}(V, A)$  where  $A$  is an alg.